SQL: Structured Query Language
Grouping Queries

SQL SELECT statement has two more clauses to support grouping operations: GROUP BY and HAVING clauses.

GROUP BY Clause

The syntax of a GROUP BY clause is

GROUP BY <AttributeName>,...,<AttributeName>

The GROUP BY clause is added to the SELECT statement after the WHERE clause (or, if there is no WHERE clause, after the FROM clause.

GROUP BY clause causes the DBMS to separate the cartesian product of all tables referenced in the FROM clause into groups of tuples.

Each group of tuples must agree on all values of attributes listed in the GROUP BY clause.

Grouping operation is used to allow for computation and reporting of aggregate operations over groups in the SELECT statement.

Consider, for example the relational table

Student(FirstName, LastName, GPA, Class, Grade, School)

The following query

SELECT School, COUNT(*)
FROM Student
GROUP BY School;

will output the number of students enrolled in each school.

The following rules need to be observed concerning the content of the SELECT clause. If a GROUP BY clause appears in the SELECT statement, then the SELECT can contain only the following:
• Attributes that are listed in the \texttt{GROUP BY} clause.
• Aggregate operations on attributes not listed in the \texttt{GROUP BY} clause.
• \texttt{COUNT(*)}

\textbf{HAVING Clause}

\texttt{GROUP BY} clause transforms the cartesian product from a relation of tuples into a relation of groups of tuples.

\texttt{HAVING} clause to the groups is what \texttt{WHERE} clause to the individual tuples. It provides a condition, and filters out the groups that fail it.

The syntax of \texttt{HAVING} clause is:

\texttt{HAVING <Condition>}

Here, the condition is a boolean combination of conditions that involve:

• Attributes from the \texttt{GROUP BY} list
• Aggregate expressions over attributes \texttt{not} from the \texttt{GROUP BY} list.

For example

\texttt{SELECT School, COUNT(*)}
\texttt{FROM Student}
\texttt{GROUP BY School}
\texttt{HAVING AVG(GPA) > 3.0;}

returns the number of students enrolled in each school, which has average student GPA over 3.0.

\textbf{Grouping and Aggregation in Relational Algebra}

Relational Algebra includes a \texttt{grouping and aggregation} operation $\gamma$.

Let $R(A_1,\ldots,A_n)$ be a relation. Let $L$ be a list $B_1,\ldots,B_k$ where $B_i$ is either:

1. one of the attributes $A_1,\ldots,A_n$ (without repetition);
2. or an expression of the form $\texttt{AGG(Exp)}[\rightarrow X]$, where $\texttt{Exp}$ is an expression over the set $\{A_1,\ldots,A_n\} - \{B_1,\ldots,B_k\}$ of attribute names, and $\texttt{AGG}$ is one of $\texttt{COUNT, MIN, MAX, SUM, AVG}$. The optional $\rightarrow X$ part, similarly to the renaming operator $\rho$ establishes an alias for the aggregate expression. As a special case, $\texttt{COUNT(*)}$ is also allowed.

Let $L = B_1,\ldots,B_k, \texttt{AGG}_1(E_1) \rightarrow X_1,\ldots,\texttt{AGG}_m(E_m) \rightarrow X_m$. The result of grouping and aggregation operation $\gamma_L(R)$ is defined below as follows:

$$
\gamma_L(R) = \{t'|\exists t \in R)((\forall i \in 1\ldots k)(t'.B_i = t.B_i)\land
(\forall j \in 1\ldots m)t'.X_j = \texttt{AGG}_j(t^*[E_j]) \text{ over the set } \{t^*\} \subset R, \text{ such that } \pi_{B_1,\ldots,B_k}(t^*) = \pi_{B_1,\ldots,B_k}(t')\}$$
Informally, the $\gamma$ operator mimics the work of the \texttt{GROUP BY} clause of the \texttt{SELECT} statement, combined with the aggregation in the \texttt{SELECT} clause. The parameter $L$ of the grouping operator represents the list of columns in the resulting relation. There are two types of columns - attributes from table $R$ - these form the groupings, and the aggregates of other columns (expressions constructed from other columns).

$$\gamma_{B_1,...,B_k,\text{AGG}_1(E_1)\rightarrow X_1,...,\text{AGG}_m(E_m)\rightarrow X_m}(R)$$

is equivalent to

\begin{verbatim}
SELECT B_1, ..., B_k, AGG_1(E_1) AS X_1, ..., AGG_m(E_m) AS X_m
FROM R
GROUP BY B_1, ..., B_k;
\end{verbatim}

To remove some of the $B_1, ..., B_k$ from the result, use \texttt{projection}.

The \texttt{HAVING} clause of the SQL \texttt{SELECT} statement is modeled via the selection operation applied to the result of the grouping and aggregation operation. In particular,

\begin{verbatim}
SELECT B_1, ..., B_k, AGG_1(E_1) AS X_1, ..., AGG_m(E_m) AS X_m
FROM R
GROUP BY B_1, ..., B_k
HAVING C;
\end{verbatim}

is equivalent to

$$\sigma_C(\gamma_{B_1,...,B_k,\text{AGG}_1(E_1)\rightarrow X_1,...,\text{AGG}_m(E_m)\rightarrow X_m}(R))$$.